Effect Handlers in Haskell, Evidently



Ningning Xie Daan Leijen

Haskell Symposium 2020 https://github.com/xnning/EvEff



monad transformers

[Liang et al. 1995]

all you need is lifting any type error can be resolved by adding more lifting

monad transformers

[Liang et al. 1995]

all you need is lifting

any type error can be resolved by adding more lifting

algebraic effects

[Plotkin and Power 2003; Plotkin and Pretnar 2013]

> composable modular

monad transformers

[Liang et al. 1995]

all you need is lifting

any type error can be resolved by adding more lifting

an alternative to

[Kammar et al. 2013; Kiselyov et al. 2013; Kiselyov and Ishii 2015; Wu and Schrijvers 2015]

algebraic effects

[Plotkin and Power 2003; Plotkin and Pretnar 2013]

composable modular

monad transformers

[Liang et al. 1995]

all you need is lifting

any type error can be resolved by adding more lifting

an alternative to

[Kammar et al. 2013; Kiselyov et al. 2013; Kiselyov and Ishii 2015; Wu and Schrijvers 2015]

[This paper]

algebraic effects

[Plotkin and Power 2003; Plotkin and Pretnar 2013]

composable modular

Contribution

ICFP 2020 Effect Handlers, Evidently

NINGNING XIE, Microsoft Research, USA JONATHAN IMMANUEL BRACHTHÄUSER, University of Tübingen, Germany DANIEL HILLERSTRÖM, The University of Edinburgh, United Kingdom PHILIPP SCHUSTER, University of Tübingen, Germany DAAN LEIJEN, Microsoft Research, USA

polymorphic algebraic effects

evidencepassing translation polymorphic evidence calculus

monadic multi-prompt translation polymorphic lambda calculus

Contribution

ICFP 2020 Effect Handlers, Evidently

NINGNING XIE, Microsoft Research, USA JONATHAN IMMANUEL BRACHTHÄUSER, University of Tübingen, Germany DANIEL HILLERSTRÖM, The University of Edinburgh, United Kingdom PHILIPP SCHUSTER, University of Tübingen, Germany DAAN LEIJEN, Microsoft Research, USA

Haskell 2020 a Haskell library Control.Ev.Eff of effect handlers via evidence-passing

polymorphic algebraic effects

evidencepassing translation oolymorphic evidence calculus

monadic multi-prompt translation polymorphic lambda calculus





Overview



- The library interface concise and simple.
- Implementations
 - 1. layered implementations
 - 2. optimizations: tail-resumptive operations are evaluated in-place
 - 3. the restriction: scoped resumptions
- Benchmarks

The library Interface

data Reader a e ans
 = Reader{ ask :: Op () a e ans }

data Reader a e ans
 = Reader{ ask :: Op () a e ans }
effect

data Reader a e ans effect
= Reader{ ask :: Op () a e ans } operation

```
data Reader a e ans
= Reader{ ask :: Op () a e ans } _____ operation
```

```
greet :: (Reader String :? e) ⇒ Eff e String
greet = do s <- perform ask ()
            return ("hello " ++ s)</pre>
```

```
data Reader a e ans
= Reader{ ask :: Op () a e ans } ---- operation
```

```
greet :: (Reader String :? e) ⇒ Eff e String
greet = do s <- perform ask ()
        return ("hello " ++ s)</pre>
```

reader action

= handler (Reader{ ask = value "world" }) action

```
effect
data Reader a e ans
  = Reader{ ask :: Op () a e ans } ---- operation
                                 effect constraint
greet :: (Reader String :? e) \Rightarrow Eff e String
                                                     effect monad
greet = do s <- perform ask ()</pre>
           return ("hello " ++ s)
reader :: Eff (Reader String :* e) ans \rightarrow Eff e ans
                                                          implementation
reader action
  = handler (Reader{ ask = value "world" }) action
```

```
effect
data Reader a e ans
  = Reader{ ask :: Op () a e ans } ---- operation
                                 effect constraint
greet :: (Reader String :? e) \Rightarrow Eff e String
                                                     effect monad
greet = do s <- perform ask ()</pre>
           return ("hello " ++ s)
reader :: Eff (Reader String :* e) ans \rightarrow Eff e ans
                                                          implementation
reader action
  = handler (Reader{ ask = value "world" }) action
```

```
data Reader a e ans
                                           effect
  = Reader{ ask :: Op () a e ans } ---- operation
                                 effect constraint
greet :: (Reader String :? e) \Rightarrow Eff e String
                                                     effect monad
greet = do s <- perform ask ()</pre>
           return ("hello " ++ s)
                                   — handles ——
reader :: Eff (Reader String :* e) ans \rightarrow Eff e ans
                                                         implementation
reader action
  = handler (Reader{ ask = value "world" }) action
```

```
data Reader a e ans
                                           effect
  = Reader{ ask :: Op () a e ans } ---- operation
                                 effect constraint
greet :: (Reader String :? e) \Rightarrow Eff e String
                                                     effect monad
greet = do s <- perform ask ()</pre>
           return ("hello " ++ s)
                                   — handles ——
reader :: Eff (Reader String :* e) ans \rightarrow Eff e ans
                                                         implementation
reader action
  = handler (Reader{ ask = value "world" }) action
```

```
data Reader a e ans
                                           effect
  = Reader{ ask :: Op () a e ans } ---- operation
                                 effect constraint
greet :: (Reader String :? e) \Rightarrow Eff e String
                                                     effect monad
greet = do s <- perform ask ()</pre>
           return ("hello " ++ s)
                                  — handles ——
reader :: Eff (Reader String :* e) ans \rightarrow Eff e ans
                                                         implementation
reader action
  = handler (Reader{ ask = value "world" }) action
```

```
helloWorld :: Eff e String
helloWorld = reader greet
```

```
data Reader a e ans
                                           effect
  = Reader{ ask :: Op () a e ans } ---- operation
                                 effect constraint
greet :: (Reader String :? e) \Rightarrow Eff e String
                                                      effect monad
greet = do s <- perform ask ()</pre>
            return ("hello " ++ s)
                                    — handles ——
reader :: Eff (Reader String :* e) ans \rightarrow Eff e ans
                                                          implementation
reader action
  = handler (Reader{ ask = value "world" }) action
 Reader String :? (Reader String :* e) 🗸
helloWorld :: Eff e String
helloWorld = reader greet
```



```
data Reader a e ans
  = Reader{ ask :: Op () a e ans }
greet :: (Reader String :? e) \Rightarrow Eff e String
greet = do s <- perform ask ()</pre>
            return ("hello " ++ s)
reader :: Eff (Reader String :* e) ans \rightarrow Eff e ans
reader action
  = handler (Reader{ ask = value "world" }) action
```

helloWorld :: Eff e String
helloWorld = reader greet

> runEff helloWorld
"hello world"

data Op a b e ans

data Op a b e ans

value :: $a \rightarrow Op$ () $a \in ans$

Reader{ ask = value "world" }

data Op a b e ans

value ::
$$a \rightarrow Op$$
 () a e ans

Reader{ ask = value "world" }

function :: $(a \rightarrow Eff e b) \rightarrow Op a b e ans$

Reader{ ask = function (\() -> "world") }

data Op a b e ans

value ::
$$a \rightarrow Op$$
 () a e ans

Reader{ ask = value "world" }

function :: $(a \rightarrow Eff e b) \rightarrow Op a b e ans$

Reader{ ask = function (\() -> "world") }

operation :: $(a \rightarrow (b \rightarrow Eff e ans) \rightarrow Eff e ans) \rightarrow Op a b e ans$

Reader{ ask = operation (\() k -> k "world") }










```
handler
  (Reader{ ask = value 1}) $
handler
  (Incr{ incr =
     operation (\x.\k. 1 + k ())}) $
handler
  (Exn{ fail =
     operation (\x.\k. 3)}) $
do x1 <- perform ask ()
     x2 <- perform ask ()
     return (x1 + x2)</pre>
```





```
handler
  (Reader{ ask = value 1}) $
handler
  (Incr{ incr =
     operation (\x.\k. 1 + k ())}) $
handler
  (Exn{ fail =
     operation (\x.\k. 3)}) $
do x1 <- perform ask ()
     x2 <- perform ask ()
     return (x1 + x2) // 2</pre>
```



```
handler
  (Reader{ ask = value 1}) $
handler
  (Incr{ incr =
     operation (\x.\k. 1 + k ())}) $
handler
  (Exn{ fail =
     operation (\x.\k. 3)}) $
do x1 <- perform ask ()
     x2 <- perform ask ()
     return (x1 + x2) // 2</pre>
```



```
handler
  (Reader{ ask = value 1}) $
handler
  (Incr{ incr =
     operation (\x.\k. 1 + k ())}) $
handler
  (Exn{ fail =
     operation (\x.\k. 3)}) $
do x1 <- perform ask ()
     x2 <- perform ask ()
     return (x1 + x2) // 2</pre>
```



```
handler
  (Reader{ ask = value 1}) $
handler
  (Incr{ incr =
     operation (\x.\k. 1 + k ())}) $
handler
  (Exn{ fail =
     operation (\x.\k. 3)}) $
do x1 <- perform ask ()
     x2 <- perform ask ()
     return (x1 + x2) // 2</pre>
```



```
handler
                                                unique marker
  (Reader{ ask = value 1}) $
                                                                       reader
handler
                                                         (m1, reader)
  (Incr{ incr =
    operation (\x.\k. 1 + k ())} $
                                                                       lincr
handler
  (Exn{ fail =
    operation (\x.\k. 3)}) $
                                                                       exception
do x1 <- perform ask ()</pre>
   x2 <- perform ask ()</pre>
   return (x1 + x2) // 2
                                                                       ask()
                                                                       ask()
```







ask()























Multi-prompt Monad

```
data Ctl a
= Pure { result :: a }
| forall ans b.
   Yield {
      marker :: Marker ans,
      op         :: (b → Ctl ans) → Ctl ans,
      cont        :: b → Ctl a }
instance Monad Ctl
```

Multi-prompt Monad

Multi-prompt Monad

```
data Ctl a
 = Pure { result :: a } ① a value result
 | forall ans b.
 Yield { ② yielding to a prompt
 the prompt to which it yields marker :: Marker ans,
 the operation impl op :: (b → Ctl ans) → Ctl ans,
the partially built up continuation cont :: b → Ctl a }
 instance Monad Ctl
```

(m3, exception) (m2, incr) (m1, reader)

(m3, exception) (m2, incr) (m1, reader)

-:: (Marker ans, Exception el ans)

(m3, exception) (m2, incr) (m1, reader)







data Context e where

```
CNil :: Context ()
CCons :: Marker ans → h e ans → Context e → Context (h :* e)
```













class h :? e where

subContext :: Context $e \rightarrow$ SubContext h










(1) instance h :? (h :* e) where subContext ctx = SubContext ctx





| 1 | instance | h | :? | (h | :* | e) | where | |
|---|----------|----|----|-----|----|-----|---------|-----|
| | subCon | te | xt | ctx | = | Sub | Context | ctx |









Overlapping instances for (h :? (h :* e))

10

1. type-level equality function

type family HEqual (h1 :: * → * → *) h2 where
HEqual h1 h1 = 'True
HEqual h1 h2 = 'False

1. type-level equality function

type family HEqual (h1 :: * → * → *) h2 where
HEqual h1 h1 = 'True
HEqual h1 h2 = 'False // datatype promotion

1. type-level equality function

type family HEqual (h1 :: * → * → *) h2 where
HEqual h1 h1 = 'True
HEqual h1 h2 = 'False

| 1. type-level | type family | HEqual (h1 :: * \rightarrow * \rightarrow *) h2 where |
|-------------------|-------------|---|
| equality function | HEqual h1 | h1 = 'True |
| | HEqual h1 | h2 = 'False |

2. effect constaints **class** (heq ~ HEqual h1 h2) \Rightarrow InEq heq h1 h2 e **where** with type equality subContextEq :: Context (h2 :* e) \rightarrow SubContext h1

1. type-leveltype family HEqual (h1 :: $* \rightarrow * \rightarrow *$) h2 whereequality functionHEqual h1 h1 = 'TrueHEqual h1 h2 = 'False

2. effect constaints **class** (heq ~ HEqual h1 h2) \Rightarrow InEq heq h1 h2 e **where** with type equality subContextEq :: Context (h2 :* e) \rightarrow SubContext h1

| 1. type-level | type family | HEqual (h1 :: * \rightarrow * \rightarrow *) h2 where |
|-------------------|-------------|---|
| equality function | HEqual h1 | h1 = 'True |
| | HEqual h1 | h2 = 'False |

2. effect constaints **class** (heq ~ HEqual h1 h2) \Rightarrow InEq heq h1 h2 e **where** with type equality subContextEq :: Context (h2 :* e) \rightarrow SubContext h1

3. delegate

| 1. type-level | type family | HEqual (h1 :: * \rightarrow * \rightarrow *) h2 where |
|-------------------|-------------|---|
| equality function | HEqual h1 | h1 = 'True |
| | HEqual h1 | h2 = 'False |

2. effect constaints **class** (heq ~ HEqual h1 h2) \Rightarrow InEq heq h1 h2 e **where** with type equality subContextEq :: Context (h2 :* e) \rightarrow SubContext h1

| 1. type-level equality function | type famil N HEqual P HEqual P | Y HEqual (h1 :: h1 h1 = 'True h1 h2 = 'False | * → * → *) h2 where | |
|---|---|---|--|--------------------|
| 2. effect constain with type equalit | ts <mark>class (hec</mark> y subConte | q ~ HEqual h1 h2 extEq :: Context |) ⇒ InEq heq h1 h2 e where (h2 : * e) → SubContext h1 | |
| (m3, h3) | (m2, h2) (m1, h1) | 1 instance | InEq 'True h1 h2 e where |) |
| | | subContextE | q ctx = SubContext ctx | |
| (m3, h3) | (m2, h2) (m1, h1) | 2 instance | | |
| | | subContextE | InEq 'False h1 h2 e wl Gq (CCons ctx) = subContext o | here ctx |
| 3. delegate | instance] subCont€ | InEq (HEqual h1 ext = subContext | h2) h1 h2 e) \Rightarrow h1 :? (h2 :* e Eq | ∋) W |

where

| 1. type-level equality function | type family HEqual (h1 :: $* \rightarrow * \rightarrow *$) h2 where HEqual h1 h1 = 'True HEqual h1 h2 = 'False |
|---|--|
| 2. effect constaints with type equality | <pre>class (heq ~ HEqual h1 h2) ⇒ InEq heq h1 h2 e where subContextEq :: Context (h2 :* e) → SubContext h1</pre> |
| (m3, h3) (m | 12, h2) (m1, h1) 1 instance InEq 'True h1 h2 e where subContextEq ctx = SubContext ctx |
| (m3, h3) (m | 12, h2) (m1, h1) ② instance InEq 'False h1 h2 e where subContextEq (CCons _ ctx) = subContext ctx |
| 3. delegate | <pre>instance InEq (HEqual h1 h2) h1 h2 e) ⇒ h1 :? (h2 :* e) where subContext = subContextEq</pre> |

| type-level equality function | <pre>type family HEqual (h1 :: * → * → *) h2 where HEqual h1 h1 = 'True HEqual h1 h2 = 'False</pre> |
|--|---|
| 2. effect constaint | S class (heq ~ HEqual h1 h2) \Rightarrow InEq heq h1 h2 e where |
| with type equality | y subContextEq :: Conterrue ~ HEqual h1 h2 text h1 |
| (m3, h3) | (m2, h2) (m1, h1) 1 instance InEq 'True h1 h2 e where |
| | <pre>subContextEq ctx = SubContext ctx</pre> |
| (m3, h3) | (m2, h2) (m1, h1) ② instance |
| | InEq 'False h1 h2 e where |
| | <pre>subContextEq (CCons ctx) = subContext ctx</pre> |
| 3. delegate | <pre>instance InEq (HEqual h1 h2) h1 h2 e) ⇒ h1 :? (h2 :* e) where subContext = subContextEq</pre> |

| 1. type-level | type family | HEqual (h1 :: * \rightarrow * \rightarrow *) h2 where |
|----------------------|--------------|---|
| equality function | HEqual h1 | h1 = 'True |
| | HEqual h1 | h2 = 'False |
| | | |
| 2. effect constaints | class (heq ~ | HEqual h1 h2) ⇒ InEq heq h1 h2 e where |
| with type equality | subContext | Eq :: Context (h2 :* e) \rightarrow SubContext h1 |

(m3, h3) (m2, h2) (m1, h1) (1) instance (h1 ~ h2) \Rightarrow InEq 'True h1 h2 e where subContextEq ctx = SubContext ctx

(m3, h3) (m2, h2) (m1, h1) ② **instance**

```
InEq 'False h1 h2 e where
```

subContextEq (CCons _ _ ctx) = subContext ctx

3. delegate

| 1. type-level | type family HEqual (h1 :: $* \rightarrow * \rightarrow *$) h2 where | |
|---------------------|--|--|
| equality function | HEqual h1 h1 = 'True | |
| | HEqual h1 h2 = 'False | |
| 2 offact constaints | | |

2. effect constaints **class** (heq ~ HEqual h1 h2) \Rightarrow InEq heq h1 h2 e **where** with type equality subContextEq :: Context (h2 :* e) \Rightarrow SubContext h1

(m3, h3)(m2, h2)(m1, h1)1instance (h1 ~ h2) \Rightarrow InEq 'True h1 h2 e where
subContextEq ctx = SubContext ctx(m3, h3)(m2, h2)(m1, h1)2instance

InEq 'False h1 h2 e where
subContextEq (CCons _ ctx) = subContext ctx

3. delegate

| 1. type-level | type family HEqual (h1 :: $* \rightarrow * \rightarrow *$) h2 where | |
|-------------------|--|--|
| equality function | HEqual h1 h1 = 'True | |
| | HEqual h1 h2 = 'False | |

2. effect constaints **class** (heq ~ HEqual h1 h2) \Rightarrow InEq heq h1 h2 e **where** with type equality subContextEq :: Context (h2 :* e) \rightarrow SubContext h1

```
(m3, h3) (m2, h2) (m1, h1) ① instance (h1 ~ h2) ⇒ InEq 'True h1 h2 e where
subContextEq ctx = SubContext ctx
(m3, h3) (m2, h2) (m1, h1) ② instance ('False ~ HEqual h1 h2
InEq 'False h1 h2 e where
subContextEq (CCons _ _ ctx) = subContext ctx
```

3. delegate

| 1. type-level | type family | HEqual (h1 :: * \rightarrow * \rightarrow *) h2 where | |
|-------------------|-------------|---|--|
| equality function | HEqual h1 | h1 = 'True | |
| | HEqual h1 | h2 = 'False | |

2. effect constaints **class** (heq ~ HEqual h1 h2) \Rightarrow InEq heq h1 h2 e **where** with type equality subContextEq :: Context (h2 :* e) \rightarrow SubContext h1

(m3, h3) (m2, h2) (m1, h1) ① instance (h1 ~ h2) ⇒ InEq 'True h1 h2 e where subContextEq ctx = SubContext ctx (m3, h3) (m2, h2) (m1, h1) ② instance ('False ~ HEqual h1 h2 , h1 :? e) ⇒ InEq 'False h1 h2 e where subContextEq (CCons _ _ ctx) = subContext ctx

3. delegate

| 1. type-level | type family | HEqual (h1 :: | * - | * | \rightarrow | *) | h2 | where |
|-------------------|-------------|---------------|-----|---|---------------|----|----|-------|
| equality function | HEqual h1 | h1 = 'True | | | | | | |
| | HEqual h1 | h2 = 'False | | | | | | |

2. effect constaints **class** (heq ~ HEqual h1 h2) \Rightarrow InEq heq h1 h2 e **where** with type equality subContextEq :: Context (h2 :* e) \rightarrow SubContext h1

```
(m3, h3) (m2, h2) (m1, h1) ① instance (h1 ~ h2) ⇒ InEq 'True h1 h2 e where
subContextEq ctx = SubContext ctx
(m3, h3) (m2, h2) (m1, h1) ② instance ('False ~ HEqual h1 h2
, h1 :? e) ⇒ InEq 'False h1 h2 e where
subContextEq (CCons _ _ ctx) = subContext ctx
```

3. delegate

















1. tail-resumptive operations (i.e., value/function) are evaluated in-place

2. non tail-resumptive operations (i.e., operation) locally decide which marker to yield to

Scoped Resumptions

- Restriction: resumptions can only be *resumed* in the same handler context as *captured*
- We believe that all important effect handlers in practice can be defined in terms of scoped resumptions
- Implemented as a dynamic check, called *guard*

- **EV** Our Control.Ev.Eff library
- **EV NT** OUR CONTROL.EV.Eff library; handlers always **N**on **T**ail-resumptive
- **EE** the Extensible Effects library [Kiselyov and Ishii 2015]
- **FE** the Fused Effects library [Schrijvers et al. 2019; Wu and Schrijvers 2015b; Wu et al. 2014]
- MTL the Monad Transformer Library

Benchmarks [Kiselyov and Ishii 2015]

| (msec) | Time Speed | | |
|--------|------------|-------|--|
| Pure | 9 | 5.44× | |
| MTL | 9 | 5.44× | |
| RunST | 41 | 1.20× | |
| EE | 339 | 0.14× | |
| FE | 10 | 4.90× | |
| EV NT | 867 | 0.06× | |
| EV | 49 | 1.00× | |

| (msec) | Time | Speed | |
|--------|------|---------------|--|
| Pure | 9 | 5.44× | |
| MTL | 9 | $5.44 \times$ | |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | $4.90 \times$ | |
| EV NT | 867 | $0.06 \times$ | |
| EV | 49 | 1.00× | |
| (msec) | Time | Speed | |
|--------|------|---------------|--|
| Pure | 9 | 5.44× | |
| MTL | 9 | $5.44 \times$ | |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | 4.90× | |
| EV NT | 867 | $0.06 \times$ | |
| EV | 49 | 1.00× | |

| (msec) | Time | Speed | |
|--------|------|---------------|--|
| Pure | 9 | 5.44× | |
| MTL | 9 | 5.44× | |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | 4.90× | |
| EV NT | 867 | 0.06× | |
| EV | 49 | 1.00× | |

| (msec) | Time | Speed | |
|--------|------|---------------|--|
| Pure | 9 | $5.44 \times$ | |
| MTL | 9 | $5.44 \times$ | |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | $4.90 \times$ | |
| EV NT | 867 | $0.06 \times$ | |
| EV | 49 | 1.00× | |

| (msec) | Time | Speed | |
|--------|------|---------------|--|
| Pure | 9 | 5.44× | |
| MTL | 9 | $5.44 \times$ | |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | $4.90 \times$ | |
| EV NT | 867 | $0.06 \times$ | |
| EV | 49 | 1.00× | |

| (msec) | Time | Speed | |
|--------|------|---------------|--|
| Pure | 9 | 5.44× | |
| MTL | 9 | $5.44 \times$ | |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | $4.90 \times$ | |
| EV NT | 867 | $0.06 \times$ | |
| EV | 49 | 1.00× | |

| (msec) | Time | Speed | |
|--------|------|---------------|---------|
| Pure | 9 | 5.44× | |
| MTL | 9 | $5.44 \times$ | |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | $4.90 \times$ | |
| EV NT | 867 | $0.06 \times$ | |
| EV | 49 | 1.00× | · |

| (msec) | Time | Speed | |
|--------|------|---------------|----------|
| Pure | 9 | 5.44× | |
| MTL | 9 | $5.44 \times$ | |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | $4.90 \times$ | — |
| EV NT | 867 | $0.06 \times$ | |
| EV | 49 | 1.00× | |

| (msec) | Time | Speed | |
|--------|------|---------------|----------|
| Pure | 9 | 5.44× | |
| MTL | 9 | $5.44 \times$ | |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | 4.90× | — |
| EV NT | 867 | 0.06× | |
| EV | 49 | 1.00× | •••••• |

```
runCount :: (State Int :? e) \Rightarrow Eff e Int
runCount = do i <- perform get ()</pre>
               if (i==0) then return i
               else do perform put (i - 1)
                        runCount
```

| $runCount_5$:: (State Integer :? e) \Rightarrow |
|--|
| Integer \rightarrow Eff e Integer |
| $runCount_5 n = foldM f 1 [n, n - 1 0]$ |
| where f acc x x 'mod' $5 == 0$ |
| = do i <- perform get () |
| perform put (i+1) |
| return (max acc x) |
| f acc x = return (max acc x) |

| (msec) | Time | Speed | |
|--------|------|---------------|----------|
| Pure | 9 | 5.44× | — |
| MTL | 9 | $5.44 \times$ | ← |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | $4.90 \times$ | — |
| EV NT | 867 | 0.06× | |
| EV | 49 | 1.00× | |

| (msec) | Time | Speed |
|--------|------|---------------|
| Pure | 247 | $0.34 \times$ |
| MTL | 327 | $0.25 \times$ |
| RunST | 256 | $0.32 \times$ |
| EE | 129 | 0.64 	imes |
| FE | 136 | $0.61 \times$ |
| EV NT | 99 | $0.84 \times$ |
| EV | 83 | 1.00 × |

```
runCount :: (State Int :? e) \Rightarrow Eff e Int
runCount = do i <- perform get ()</pre>
               if (i==0) then return i
               else do perform put (i - 1)
                        runCount
```

| runCount ₅ :: (State Integer :? e) \Rightarrow |
|---|
| Integer \rightarrow Eff e Integer |
| runCount ₅ n = foldM f 1 [n, n - 1 0] |
| where f acc x x 'mod' $5 == 0$ |
| = do i <- perform get () |
| perform put (i+1) |
| return (max acc x) |
| f acc x = return (max acc x) |

| (msec) | Time | Speed | |
|--------|------|---------------|----------|
| Pure | 9 | 5.44 	imes | |
| MTL | 9 | 5.44 	imes | — |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | 0.14 	imes | |
| FE | 10 | $4.90 \times$ | |
| EV NT | 867 | $0.06 \times$ | |
| EV | 49 | 1.00 × | |

| (msec) | Time | Speed | |
|--------|------|---------------|--|
| Pure | 247 | $0.34 \times$ | |
| MTL | 327 | $0.25 \times$ | |
| RunST | 256 | $0.32 \times$ | |
| EE | 129 | $0.64 \times$ | |
| FE | 136 | $0.61 \times$ | |
| EV NT | 99 | $0.84 \times$ | |
| EV | 83 | 1.00× | |

```
runCount :: (State Int :? e) \Rightarrow Eff e Int
runCount = do i <- perform get ()</pre>
               if (i==0) then return i
               else do perform put (i - 1)
                        runCount
```

| runCount ₅ :: (State Integer :? e) \Rightarrow |
|---|
| Integer \rightarrow Eff e Integer |
| $runCount_5 n = foldM f 1 [n, n - 1 0]$ |
| where f acc $x \mid x \pmod{5} = 0$ |
| = do i <- perform get () |
| perform put (i+1) |
| return (max acc x) |
| f acc x = return (max acc x) |

| (msec) | Time | Speed | |
|--------|------|---------------|---|
| Pure | 9 | $5.44 \times$ | - |
| MTL | 9 | $5.44 \times$ | - |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | $4.90 \times$ | - |
| EV NT | 867 | $0.06 \times$ | |
| EV | 49 | 1.00 × | |

| (msec) | Time | Speed | |
|--------|------|---------------|--|
| Pure | 247 | 0.34× | |
| MTL | 327 | $0.25 \times$ | |
| RunST | 256 | 0.32× | |
| EE | 129 | 0.64× | |
| FE | 136 | 0.61× | |
| EV NT | 99 | $0.84 \times$ | |
| EV | 83 | 1.00× | |

| runCount ₅ :: (State Integer :? e) \Rightarrow |
|---|
| Integer \rightarrow Eff e Integer |
| $runCount_5 n = foldM f 1 [n, n - 1 0]$ |
| where f acc x x 'mod' $5 == 0$ |
| = do i <- perform get () |
| perform put (i+1) |
| return (max acc x) |
| f acc x = return (max acc x) |

| (m | sec) | Time | Speed | | |
|-----|--------------|------|---------------|----------|----|
| Pι | ıre | 247 | $0.34 \times$ | 4 | |
| M | TL | 327 | $0.25 \times$ | (| |
| Rı | ınST | 256 | $0.32 \times$ | (| |
| EF | <u>-</u> | 129 | 0.64 	imes | | I |
| FE | 5 | 136 | $0.61 \times$ | + | I |
| E E | V NT | 99 | $0.84 \times$ | | |
| E | \checkmark | 83 | 1.00× | | 16 |

| (msec) | Time | Speed | |
|--------|------|---------------|----------|
| Pure | 9 | $5.44 \times$ | |
| MTL | 9 | $5.44 \times$ | — |
| RunST | 41 | $1.20 \times$ | |
| EE | 339 | $0.14 \times$ | |
| FE | 10 | $4.90 \times$ | |
| EV NT | 867 | 0.06× | |
| EV | 49 | 1.00× | |











In the benchmark, we put many **Reader** layers under or over the target **State** layer.



In the benchmark, we put many **Reader** layers under or over the target **State** layer.



In the benchmark, we put many **Reader** layers under or over the target **State** layer.









18

More in the Paper

• Advanced handlers:

handlers with return clauses, handlers with local state

- Our implementation ensures type safety
- More discussion





More in the Paper

• Advanced handlers:

handlers with return clauses, handlers with local state

- Our implementation ensures type safety
- More discussion

Effect Handlers, Evidently

Ningning Xie Jonathan Brachthäuser Daniel Hillerström Philipp Schuster Daan Leijen ICFP 2020



https://github.com/xnning/EvEff